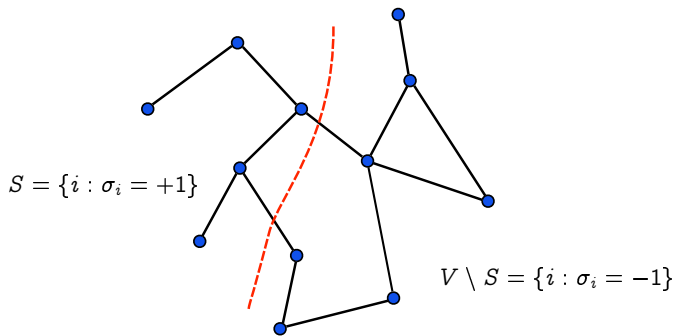


Extremal cuts:
From sparse random graphs to spin-glasses

Amir Dembo

July 20, 2016

Cuts in graph $G = (V, E)$



$$\text{cut}_G(S) \equiv \left| \{(i, j) \in E : i \in S, j \in V \setminus S\} \right|$$

$$2\text{cut}_G(S) - |E| = - \sum_{(i,j) \in E} \sigma_i \sigma_j \quad \text{AF-Ising on } G$$

Extremal cuts

Minimum bisection

$$\text{mcut}(G) = \min \left\{ \text{cut}_G(S) : S \subseteq V, |S| = |V|/2 \right\}$$

Maximum bisection

$$\text{MCUT}(G) = \max \left\{ \text{cut}_G(S) : S \subseteq V, |S| = |V|/2 \right\}$$

Maximum Cut

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Complexity: hard to solve (worst case)

mcut
MCUT
MaxCut

} are NP-hard

(hard to approximate to within $1 + o(1)$ factor:

SDP [Goemans,Williamson '95] gives 0.878

hardness [Trevisan,Sorkin,Sudan,Williamson '00] best be 0.941)

Typical complexity (average graph)?

Random graph models

Erdős-Renyi Random graph $G = (V, E) \sim \mathcal{G}(n, p)$

- ▶ $|V| = n$ vertices
- ▶ Edges independently present with probability p
(or uniform among graphs of $m = \binom{n}{2}p$ edges)
- ▶ Average degree $\gamma = np$

Random regular graph $G = (V, E) \sim \mathcal{G}^{\text{reg}}(n, \gamma)$

- ▶ $|V| = n$ vertices
- ▶ Uniformly random among graphs with $\deg(i) = \gamma \quad \forall i \in V$.

$p = \gamma/n, \gamma = O(1)$, sparse graphs

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Extremal cuts of random graphs: long history

- ▶ Bollobas '88 (\mathcal{G}^{reg} , mcut, concentration bound)
- ▶ Alon '97 (\mathcal{G}^{reg} , mcut, algorithmic bound)
- ▶ Coppersmith, Gamarnik, Hajiaghayi, Sorkin '04 (\mathcal{G} , MaxCut transition $\gamma = 0.5$)
- ▶ Díaz, Serna, Wormald, '07 (\mathcal{G}^{reg} , MCUT, algorithmic bound)
- ▶ Daudé, Martínez, Rasendrasahina, Ravelomanana '12 (MaxCut, scaling-window)
- ▶ Gamarnik, Li '14 (\mathcal{G} , MaxCut, improved concentration)
- ▶ ...

Typical result (sparse case):

If $G \sim \mathcal{G}(n, \gamma/n)$ then, with high probability

$$\frac{n\gamma}{4} + C_1 n \sqrt{\gamma} \leq \text{MCUT}(G) \leq \frac{n\gamma}{4} + C_2 n \sqrt{\gamma}$$

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A classical argument

(Bollobas '88)

$|V| = n, |E| = n\gamma/2.$ Fix $S \subseteq V, |S| = n/2$

- ▶ Each edge is cut with probability $1/2$

$$\mathbb{E}[\text{cut}(S)] = \frac{|E|}{2} = \frac{n\gamma}{4}.$$

(also factor 0.5 approximation of MaxCut)

- ▶ Azuma-Hoeffding argument

$$\mathbb{P}\left\{|\text{cut}(S) - \mathbb{E}\text{cut}(S)| \geq \Delta\right\} \leq 2 \exp\left(-\frac{\Delta^2}{4n\gamma}\right)$$

(simpler: Binomial($\frac{n^2}{4}, \frac{\gamma}{n}$) in $\mathcal{G}(n, \gamma/n)$)

- ▶ Union bound

$$\mathbb{P}\left\{\max_{S, |S|=n/2} \left|\text{cut}(S) - \frac{n\gamma}{4}\right| \geq \delta n \sqrt{\gamma}\right\} \leq 2^{n+1} e^{-n\delta^2/4}.$$

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$$\frac{n\gamma}{4} - C_1 n\sqrt{\gamma} \leq \text{mcut}(G) \leq \text{MCUT}(G) \leq \frac{n\gamma}{4} + C_2 n\sqrt{\gamma}$$

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Extremal cuts and Ising spins

Ising Hamiltonian: $\mathcal{H}_G(\sigma) = \frac{1}{\sqrt{\gamma}} \sum_{(i,j) \in E} \sigma_i \sigma_j : \{-1, +1\}^n \mapsto \mathbb{R}.$

$$\begin{aligned} \text{MaxCut}(G) &= \max_{\sigma} \sum_{(i,j) \in E} \left(\frac{1 - \sigma_i \sigma_j}{2} \right) \\ &= \frac{1}{2} |E| - \frac{1}{2} \min_{\sigma} \sum_{(i,j) \in E} \sigma_i \sigma_j = \frac{n\gamma}{4} - \sqrt{\frac{\gamma}{4}} \min_{\sigma} \{\mathcal{H}_G(\sigma)\} \end{aligned}$$

$$(S = \{i : \sigma_i = +1\}, \quad |E| = \frac{n\gamma}{2}).$$

$$|S| = \frac{n}{2} \iff \Omega_n = \left\{ \sigma \in \{-1, +1\}^n : \sum_{i=1}^n \sigma_i = 0 \right\}$$

$$\text{MCUT}(G) = \frac{n\gamma}{4} - \sqrt{\frac{\gamma}{4}} \min_{\sigma \in \Omega_n} \{\mathcal{H}_G(\sigma)\},$$

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Statistical physics: Ising measures

Ising measures (on balanced configurations):

$$\mathcal{H}_G(\sigma) = \frac{1}{\sqrt{\gamma}} \sum_{(i,j) \in E} \sigma_i \sigma_j \quad (\sigma \in \Omega_n),$$
$$\mu_{n,\beta}(\sigma) \equiv \frac{1}{Z_n(\beta)} e^{\beta \mathcal{H}_G(\sigma)}, \quad Z_n(\beta) \equiv \sum_{\tilde{\sigma} \in \Omega_n} e^{\beta \mathcal{H}_G(\tilde{\sigma})}.$$

(ferromagnetic iff $\beta \geq 0$).

Ground state energy:

$$\lim_{\beta \rightarrow +\infty} \frac{1}{\beta} \log Z_n(\beta) = \max_{\sigma \in \Omega_n} \{ \mathcal{H}_G(\sigma) \} \quad \Longrightarrow \quad \text{mcut}(G),$$
$$\lim_{\beta \rightarrow -\infty} \frac{1}{\beta} \log Z_n(\beta) = \min_{\sigma \in \Omega_n} \{ \mathcal{H}_G(\sigma) \} \quad \Longrightarrow \quad \text{MCUT}(G).$$

Insights from statistical physics

Conjectures

- ▶ Fu, Anderson, '86
- ▶ Mézard, Parisi, '01
- ▶ [Zdéborova, Boettcher '10] For $G_n \sim \mathcal{G}^{\text{reg}}(n, \gamma)$, w.h.p.

$$\text{MCUT}(G_n) = \text{MaxCut}(G_n) + o(n) = |E_n| - \text{mcut}(G_n) + o(n).$$

Theorem (Bayati, Gamarnik, Tetali, '09)

For $G_n \sim \mathcal{G}(n, \gamma/n)$, or $G_n \in \mathcal{G}^{\text{reg}}(n, \gamma)$ w.h.p.

$$\liminf_{n \rightarrow \infty} \frac{1}{n} \text{MaxCut}(G_n) = \limsup_{n \rightarrow \infty} \frac{1}{n} \text{MaxCut}(G_n)$$

Sub-additivity: no limit value.

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Connection to spin glasses

Theorem (D.-Montanari-Sen '15)

For $G_n \sim \mathcal{G}^{\text{reg}}(n, \gamma)$, $G_n \sim \mathcal{G}(n, \gamma/n)$ w.h.p.

$$\begin{aligned}\frac{1}{n} \text{mcut}(G_n) &= \frac{\gamma}{4} - P_* \sqrt{\frac{\gamma}{4}} - o_\gamma(\sqrt{\gamma}) + o_n(1), \\ \frac{1}{n} \text{MCUT}(G_n) &= \frac{\gamma}{4} + P_* \sqrt{\frac{\gamma}{4}} + o_\gamma(\sqrt{\gamma}) + o_n(1), \\ \frac{1}{n} \text{MaxCut}(G_n) &= \frac{\gamma}{4} + P_* \sqrt{\frac{\gamma}{4}} + o_\gamma(\sqrt{\gamma}) + o_n(1).\end{aligned}$$

What is P_* ?

GOE random matrix:

$$J \in \mathbb{R}^{n \times n}, \quad J = J^T, \quad (J_{ij})_{i < j} \sim N(0, n^{-1}), \quad J_{ii} \sim N(0, 2n^{-1}).$$

Sherrington-Kirkpatrick spin-glass model

$$\mathcal{H}_{\text{SK}}(\sigma) \equiv \frac{1}{2} \sigma^T J \sigma.$$

Finally

$$P_* \equiv \lim_{n \rightarrow \infty} \frac{1}{n} \mathbb{E} \left[\max_{\sigma \in \{-1, +1\}^n} \{\mathcal{H}_{\text{SK}}(\sigma)\} \right].$$

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- ▶ The limit exists (consequence of [Guerra, Toninelli '02])
- ▶ Given by '*Parisi's formula*' ([Talagrand '06])
- ▶ Clarifies why 'standard' combinatorial methods unsuccessful.
- ▶ Set of near-extremal cuts **predicted** to have ∞ -RSB structure.

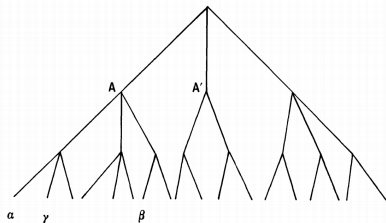


Fig. 1. — The tree of the states. The different states $\alpha, \beta, \gamma, \dots$ are the extremities of the branches of the tree. The distance between two states is a monotonic function of the number of steps one has to climb along the tree to find a common ancestor.

Proof strategy

1. Proof for mcut, MCUT with $G \sim \mathcal{G}(n, \gamma/n)$:

Concentration, interpolation, 'smooth max'

2. Extend to $G' \sim \mathcal{G}^{\text{reg}}(n, \gamma)$: Couple G' and $G \sim \mathcal{G}(n, \gamma'/n)$

(Tricky: need $\gamma - \gamma' \gg \sqrt{\gamma}$ so $|E_G \Delta E_{G'}| \gg n\sqrt{\gamma}$)

3. Prove that $\text{MaxCut}(G) - \text{MCUT}(G) = o(n\sqrt{\gamma})$

(Tricky part: $\frac{n}{2} - |S| \leq n\gamma^{-\delta}$, $\delta < \frac{1}{4}$)

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mcut for $\mathcal{G}(n, \gamma/n)$:

Recall (our claim)

$$\text{mcut}(G) = \frac{n\gamma}{4} - \sqrt{\frac{\gamma}{4}} \max_{\sigma \in \Omega_n} \{\mathcal{H}_G(\sigma)\},$$

By concentration, suffices to show

$$\frac{1}{n} \mathbb{E} \left[\max_{\sigma \in \Omega_n} \{\mathcal{H}_G(\sigma)\} \right] = P_* + o_\gamma(1) + o_n(1).$$

We know

$$\frac{1}{n} \mathbb{E} \left[\max_{\sigma \in \{-1, +1\}^n} \{\mathcal{H}_{\text{SK}}(\sigma)\} \right] = P_* + o_n(1).$$

Idea: Interpolation ('smart path'; Lindeberg method)

Interpolation: 'smart path'

Step 1: For $t \in [0, 1]$, $G(t) = (V, E(t)) \sim \mathcal{G}(n, \gamma(1-t)/n)$

$$\mathcal{H}_t(\sigma; \gamma) = \frac{1}{\sqrt{\gamma}} \sum_{(i,j) \in E(t)} \sigma_i \sigma_j + \frac{\sqrt{t}}{2} \sum_{i,j=1}^n J_{ij} \sigma_i \sigma_j$$

$$\mathcal{H}_0 = \mathcal{H}_G \quad \mathcal{H}_1 = \mathcal{H}_{\text{SK}} \quad (\sigma \in \Omega_n).$$

$G(t)$ multi-graph of i.i.d. Poisson($\gamma(1-t)/n$) multi-edges.

Step 2: From $\max_{\sigma \in \Omega_n} \{\mathcal{H}_{\text{SK}}(\sigma)\}$ to $\max_{\sigma \in \{-1, +1\}^n} \{\mathcal{H}_{\text{SK}}(\sigma)\}$.

Step 1: 'smooth max'

Free energy density (balanced configurations):

$$\phi_n(\beta; t, \gamma) \equiv \frac{1}{n} \mathbb{E} \left[\log \left\{ \sum_{\sigma \in \Omega_n} e^{\beta \mathcal{H}_t(\sigma; \gamma)} \right\} \right]$$
$$\left| \frac{1}{\beta} \phi_n(\beta; t, \gamma) - \frac{1}{n} \mathbb{E} \left[\max_{\sigma \in \Omega_n} \{ \mathcal{H}_t(\sigma; \gamma) \} \right] \right| \leq \frac{\log 2}{\beta}.$$

\implies just find $\beta(\gamma) \rightarrow \infty$ with

$$\sup_{n, t} \frac{1}{\beta} \left| \frac{\partial \phi_n}{\partial t}(\beta; t, \gamma) \right| \rightarrow 0.$$

Step 1: Poisson & Gaussian calculus

Lemma

$$\left| \frac{\partial \phi_n}{\partial t}(\beta; t, \gamma) \right| \leq C \left[\frac{\beta^3}{\gamma^{1/2}} + \frac{\beta^4}{\gamma} \right].$$

Proof:

$\langle \cdot \rangle_{n,t}$ denote expectations for i.i.d. $\sigma^{(j)} \in \Omega_n$ under $\mu_{n,t,\beta,J}(\sigma) \equiv Z_{n,t,J}^{-1} e^{\beta \mathcal{H}_t(\sigma; \gamma)}$.

Multi-replica overlaps $Q_\ell \equiv \frac{1}{n} \sum_{i=1}^n \left(\prod_{j=1}^{\ell} \sigma_i^{(j)} \right),$

so $Q_1 = 0$ and $|Q_\ell| \leq 1$. Let $\kappa \equiv \beta / \sqrt{\gamma}$.

$$\frac{\partial \phi_n(\beta; t, \gamma)}{\partial t} = I + II,$$

$$II = \frac{\beta^2}{4} (1 - \mathbb{E}[\langle Q_2^2 \rangle_{n,t}]) \quad \text{by Gaussian (SK) calculus}$$

$$I = \frac{\gamma}{2} \left(-\log \cosh \kappa + \sum_{\ell \geq 1} \frac{(-\tanh \kappa)^\ell}{\ell} \mathbb{E}[\langle Q_\ell^2 \rangle_{n,t}] \right) \quad \text{by Poisson (ER) calculus}$$

Step 1: Poisson & Gaussian calculus

Lemma

$$\left| \frac{\partial \phi_n}{\partial t}(\beta; t, \gamma) \right| \leq C \left[\frac{\beta^3}{\gamma^{1/2}} + \frac{\beta^4}{\gamma} \right].$$

Proof:

$\langle \cdot \rangle_{n,t}$ denote expectations for i.i.d. $\sigma^{(j)} \in \Omega_n$ under $\mu_{n,t,\beta,J}(\sigma) \equiv Z_{n,t,J}^{-1} e^{\beta \mathcal{H}_t(\sigma; \gamma)}$.

Multi-replica overlaps $Q_\ell \equiv \frac{1}{n} \sum_{i=1}^n \left(\prod_{j=1}^{\ell} \sigma_i^{(j)} \right),$

so $Q_1 = 0$ and $|Q_\ell| \leq 1$. Let $\kappa \equiv \beta / \sqrt{\gamma}$.

$$\frac{\partial \phi_n(\beta; t, \gamma)}{\partial t} = I + II,$$

$$II = \frac{\beta^2}{4} (1 - \mathbb{E}[\langle Q_2^2 \rangle_{n,t}]) \quad \text{by Gaussian (SK) calculus}$$

$$I = \frac{\gamma}{2} \left(-\log \cosh \kappa + \sum_{\ell \geq 1} \frac{(-\tanh \kappa)^\ell}{\ell} \mathbb{E}[\langle Q_\ell^2 \rangle_{n,t}] \right) \quad \text{by Poisson (ER) calculus}$$

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Universality [Subhabrata Sen, '16]

Not Ising related, for all models when $\gamma \gg 1$ (& ER).

Key: replace 'smart path' with Lindeberg method.

Applications:

- ▶ Max q -cut in sparse random graphs

\implies Potts spin-glass.

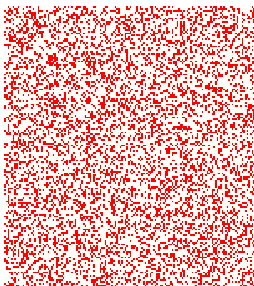
- ▶ Max-SAT for $m = \gamma n/2$ random linear Eqn. over $\{0, 1\}$
(in UNSAT phase, k variables per Eqn.)

\implies k -spin models.

Finding $\text{mcut}(G)$ for specific $G \sim \mathcal{G}(n, \gamma/n)$

$$\text{mcut}(G) \iff \max \langle \sigma, A_G \sigma \rangle,$$

for $\sum_{i=1}^n \sigma_i = 0, \quad \sigma_i \in \{-1, 1\}.$



$$(A_G)_{ij} = \begin{cases} 1 & \text{if } (i, j) \in E, \\ 0 & \text{otherwise.} \end{cases}$$

Regularization (Lagrangian)

$$\begin{aligned} \max \quad & \langle \sigma, A_G \sigma \rangle - \lambda \left(\sum_{i=1}^n \sigma_i \right)^2, \\ \text{for } & \sigma_i \in \{-1, 1\}. \end{aligned}$$

$G \sim \mathcal{G}(n, \gamma/n) \Rightarrow$ 'sufficient' regularization $\lambda = \gamma/n$:

$$\begin{aligned} \text{OPT}(G) = \max \quad & \langle \sigma, (A_G - \mathbb{E}\{A_G\})\sigma \rangle, \\ \text{for } & \sigma_i \in \{-1, 1\}. \end{aligned}$$

Recall for such G w.h.p.

$$\text{OPT}(G) = 2nP_*\sqrt{\gamma} + n o_\gamma(\sqrt{\gamma}) \approx 1.5264 n\sqrt{\gamma}$$

Spectral relaxation: does poorly!

$$\text{OPT}(G) = \max \langle \sigma, (A_G - \mathbb{E}\{A_G\})\sigma \rangle, \\ \text{for } \sigma_i \in \{+1, -1\} .$$

$$\lambda_{\max}(A_G - \mathbb{E}A_G) n = \max \langle \sigma, (A_G - \mathbb{E}\{A_G\})\sigma \rangle, \\ \text{for } \sigma_i \in \{\cancel{+1}, -1\} \quad \|\sigma\|_2^2 = n$$

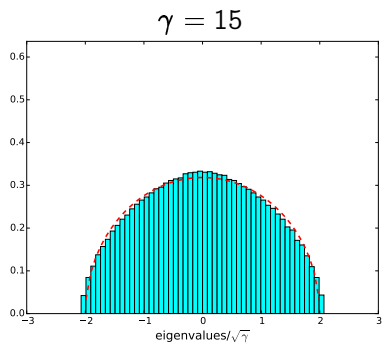
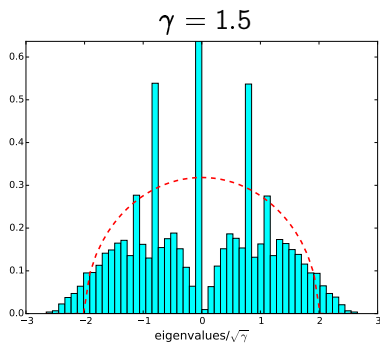
Theorem (Krivelevich, Sudakov 2003+Vu 2005)

If $G \sim \mathcal{G}(n, \gamma/n)$, then, with high probability,

$$\lambda_{\max}(A_G - \mathbb{E}A_G) = \begin{cases} 2\sqrt{\gamma}(1 + o(1)) & \text{if } \gamma \gg (\log n)^4, \\ \sqrt{\log n / (\log \log n)}(1 + o(1)) & \text{if } \gamma = O(1). \end{cases}$$

Compare with $\text{OPT}(G)/n \approx 1.5264 \sqrt{\gamma}$

Illustration: $n = 10,000$



Semidefinite Programming - does well!

$$\text{OPT}(G) = \max \langle (A_G - \mathbb{E}\{A_G\}), \sigma\sigma^\top \rangle, \\ \text{for } \sigma_i \in \{+1, -1\}.$$

$$\text{SDP}(G) = \max \langle (A_G - \mathbb{E}\{A_G\}), X \rangle, \\ \text{for } X_{ii} = 1, X \succeq 0.$$

Theorem (Montanari-Sen '15)

If $G \sim \mathcal{G}(n; \gamma/n)$, $\gamma = O(1)$, *w.h.p.*

$$\frac{1}{n} \text{SDP}(G) = 2\sqrt{\gamma} + o(\sqrt{\gamma})$$

- ▶ $\text{SDP}(G)$ behaves much better than the principal eigenvalue

Semidefinite Programming - does well!

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Proof outline

'Interpolate' in rank:

$$\begin{aligned} \text{OPT}_k(G) &:= \max \langle (A_G - \mathbb{E}\{A_G\}), X \rangle, \\ &\text{for } X_{ii} = 1, X \succeq 0, \text{rank}(X) \leq k \end{aligned}$$

$$\text{OPT} = \text{OPT}_1 \nearrow \text{OPT}_n = \text{SDP}.$$

Equivalently:

$$\begin{aligned} \text{OPT}_k(G) &= \max \langle (A_G - \mathbb{E}\{A_G\}), \sigma\sigma^\top \rangle, \\ &\text{for } \sigma = (\sigma_1, \sigma_2, \dots, \sigma_n)^\top \in \mathbb{R}^{n \times k}, \\ &\sigma_i \in \mathbb{R}^k, \|\sigma_i\|_2 = 1. \end{aligned}$$

$$\begin{aligned} \text{OPT}_k(G) &= \max \sum_{i,j=1}^n (A_G - \mathbb{E}\{A_G\})_{i,j} \langle \sigma_i, \sigma_j \rangle, \\ &\text{for } \sigma_i \in \mathbb{R}^k, \|\sigma_i\|_2 = 1. \end{aligned}$$

\mathbb{S}^{k-1} -valued spin model

Proof outline

- ▶ Higher-rank Grothendieck inequality

$$\left(1 - \frac{c}{k}\right) \text{SDP}(G) \leq \text{OPT}_k(G) \leq \text{SDP}(G).$$

- ▶ Interpolation for $\text{OPT}_k(G)$ (via asymptotic free energy)

$$(A_G - \mathbb{E}\{A_G\}) \longleftrightarrow J \sim \text{GOE}_n$$

- ▶ Analyze SDP for $J \sim \text{GOE}_n$.

Conclusion

- ▶ Extremal cuts in random graphs
 \implies balanced Ising spins $\sigma_i \in \{-1, +1\}$
- ▶ SDP \implies Vector spins $\sigma_i \in \mathbb{S}^{k-1}$, $k \leq n$.
- ▶ Universal connection to spin-glasses as $\gamma \rightarrow \infty$.
- ▶ But little is known for γ fixed!

Open Problem [Zdéborova, Boettcher '10]:

For $G_n \sim \mathcal{G}^{\text{reg}}(n, \gamma)$, **fixed** $\gamma \geq 3$, w.h.p.

$$\text{MCUT}(G_n) = \text{MaxCut}(G_n) + o(n) = |E_n| - \text{mcut}(G_n) + o(n).$$