

Fast Fourier Transform: Why? How?

Shamgar Gurevich

UW Madison

May 13, 2012

Joint work with:

Ronny Hadani (Math, Austin)

Motivation:

- Discrete Fourier Transform

$$DFT = \frac{1}{\sqrt{N}} \left(e^{\frac{2\pi i}{N} \tau \cdot \omega} \right)_{0 \leq \tau, \omega \leq N-1}$$

Motivation:

- Discrete Fourier Transform

$$DFT = \frac{1}{\sqrt{N}} \left(e^{\frac{2\pi i}{N} \tau \cdot \omega} \right)_{0 \leq \tau, \omega \leq N-1}$$

- Compute:

$$\hat{f} = DFT[f]; \quad f = \begin{pmatrix} f(0) \\ \cdot \\ \cdot \\ \cdot \\ f(N-1) \end{pmatrix}; \quad \text{Fast!!}$$

Motivation:

- Discrete Fourier Transform

$$DFT = \frac{1}{\sqrt{N}} \left(e^{\frac{2\pi i}{N} \tau \cdot \omega} \right)_{0 \leq \tau, \omega \leq N-1}$$

- Compute:

$$\hat{f} = DFT[f]; \quad f = \begin{pmatrix} f(0) \\ \vdots \\ f(N-1) \end{pmatrix}; \quad \text{Fast!!}$$

- Cooley–Tukey (1965): $O(N \cdot \log(N))$ operations!

Solution:

(I) Heisenberg Group Representation

- $\mathcal{H} = \mathbb{C}(\mathbb{Z}_N)$ — Hilbert space of digital signals.

Solution:

(I) Heisenberg Group Representation

- $\mathcal{H} = \mathbb{C}(\mathbb{Z}_N)$ — Hilbert space of digital signals.
 - $f : \{0, \dots, N - 1\} \rightarrow \mathbb{C}$.

Solution:

(I) Heisenberg Group Representation

- $\mathcal{H} = \mathbb{C}(\mathbb{Z}_N)$ — Hilbert space of digital signals.
 - $f : \{0, \dots, N - 1\} \rightarrow \mathbb{C}$.
- Basic operations

Solution:

(I) Heisenberg Group Representation

- $\mathcal{H} = \mathbb{C}(\mathbb{Z}_N)$ — Hilbert space of digital signals.
 - $f : \{0, \dots, N-1\} \rightarrow \mathbb{C}$.
- Basic operations
 - Time shift: $\tau \in \mathbb{Z}_N$,

$$L_\tau : \mathcal{H} \rightarrow \mathcal{H},$$
$$L_\tau[f](t) = f(t + \tau), \quad t \in \mathbb{Z}_N.$$

Solution:

(I) Heisenberg Group Representation

- $\mathcal{H} = \mathbb{C}(\mathbb{Z}_N)$ — Hilbert space of digital signals.
 - $f : \{0, \dots, N-1\} \rightarrow \mathbb{C}$.

- Basic operations

- Time shift: $\tau \in \mathbb{Z}_N$,

$$L_\tau : \mathcal{H} \rightarrow \mathcal{H},$$
$$L_\tau[f](t) = f(t + \tau), \quad t \in \mathbb{Z}_N.$$

- Frequency shift: $\omega \in \mathbb{Z}_N$,

$$M_\omega : \mathcal{H} \rightarrow \mathcal{H},$$
$$M_\omega[f](t) = e^{\frac{2\pi i}{N}\omega t} f(t).$$

Solution:

(I) Heisenberg Group Representation

- $\mathcal{H} = \mathbb{C}(\mathbb{Z}_N)$ — Hilbert space of digital signals.
 - $f : \{0, \dots, N-1\} \rightarrow \mathbb{C}$.

- Basic operations

- Time shift: $\tau \in \mathbb{Z}_N$,

$$L_\tau : \mathcal{H} \rightarrow \mathcal{H},$$
$$L_\tau[f](t) = f(t + \tau), \quad t \in \mathbb{Z}_N.$$

- Frequency shift: $\omega \in \mathbb{Z}_N$,

$$M_\omega : \mathcal{H} \rightarrow \mathcal{H},$$
$$M_\omega[f](t) = e^{\frac{2\pi i}{N}\omega t} f(t).$$

- Note:

$$M_\omega \circ L_\tau = e^{\frac{2\pi i}{N}\omega\tau} \cdot L_\tau \circ M_\omega \quad \text{— Heisenberg commutation relations}$$

Heisenberg Rep'n, cont.

- Combine: $\tau, \omega, z \in \mathbb{Z}_N$

$$\pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau.$$

Heisenberg Rep'n, cont.

- Combine: $\tau, \omega, z \in \mathbb{Z}_N$

$$\pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau.$$

- Identity:

$$\pi(\tau, \omega, z) \circ \pi(\tau', \omega', z') = \pi(\tau + \tau', \omega + \omega', z + z' + \frac{1}{2} \begin{vmatrix} \tau & \omega \\ \tau' & \omega' \end{vmatrix}).$$

Heisenberg Rep'n, cont.

- Combine: $\tau, \omega, z \in \mathbb{Z}_N$

$$\pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau.$$

- Identity:

$$\pi(\tau, \omega, z) \circ \pi(\tau', \omega', z') = \pi\left(\tau + \tau', \omega + \omega', z + z' + \frac{1}{2} \begin{vmatrix} \tau & \omega \\ \tau' & \omega' \end{vmatrix}\right).$$

- Question. How to think on this?

Heisenberg Rep'n, cont.

- Combine: $\tau, \omega, z \in \mathbb{Z}_N$

$$\pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau.$$

- Identity:

$$\pi(\tau, \omega, z) \circ \pi(\tau', \omega', z') = \pi\left(\tau + \tau', \omega + \omega', z + z' + \frac{1}{2} \begin{vmatrix} \tau & \omega \\ \tau' & \omega' \end{vmatrix}\right).$$

- Question. How to think on this?
- Answer. Heisenberg group:

Heisenberg Rep'n, cont.

- Combine: $\tau, \omega, z \in \mathbb{Z}_N$

$$\pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau.$$

- Identity:

$$\pi(\tau, \omega, z) \circ \pi(\tau', \omega', z') = \pi(\tau + \tau', \omega + \omega', z + z' + \frac{1}{2} \begin{vmatrix} \tau & \omega \\ \tau' & \omega' \end{vmatrix}).$$

- Question. How to think on this?
- Answer. Heisenberg group:

- $H = \underbrace{\mathbb{Z}_N \times \mathbb{Z}_N}_V \times \underbrace{\mathbb{Z}_N}_Z;$

Heisenberg Rep'n, cont.

- Combine: $\tau, \omega, z \in \mathbb{Z}_N$

$$\pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau.$$

- Identity:

$$\pi(\tau, \omega, z) \circ \pi(\tau', \omega', z') = \pi\left(\tau + \tau', \omega + \omega', z + z' + \frac{1}{2} \begin{vmatrix} \tau & \omega \\ \tau' & \omega' \end{vmatrix}\right).$$

- Question. How to think on this?

- Answer. Heisenberg group:

- $H = \underbrace{\mathbb{Z}_N \times \mathbb{Z}_N}_v \times \underbrace{\mathbb{Z}_N}_z;$

- $(v, z) \cdot (v', z') = \left(v + v', z + z' + \frac{1}{2} \begin{vmatrix} v & \\ v' & \end{vmatrix} \right);$

Heisenberg Rep'n, cont.

- Combine: $\tau, \omega, z \in \mathbb{Z}_N$

$$\pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau.$$

- Identity:

$$\pi(\tau, \omega, z) \circ \pi(\tau', \omega', z') = \pi\left(\tau + \tau', \omega + \omega', z + z' + \frac{1}{2} \begin{vmatrix} \tau & \omega \\ \tau' & \omega' \end{vmatrix}\right).$$

- Question. How to think on this?

- Answer. Heisenberg group:

- $H = \underbrace{\mathbb{Z}_N \times \mathbb{Z}_N}_v \times \underbrace{\mathbb{Z}_N}_z;$

- $(v, z) \cdot (v', z') = (v + v', z + z' + \frac{1}{2} \begin{vmatrix} v & z \\ v' & z' \end{vmatrix});$

- $(0, 0) \cdot (v, z) = (v, z) \cdot (0, 0) = (v, z);$

Heisenberg Rep'n, cont.

- Combine: $\tau, \omega, z \in \mathbb{Z}_N$

$$\pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau.$$

- Identity:

$$\pi(\tau, \omega, z) \circ \pi(\tau', \omega', z') = \pi(\tau + \tau', \omega + \omega', z + z' + \frac{1}{2} \begin{vmatrix} \tau & \omega \\ \tau' & \omega' \end{vmatrix}).$$

- Question. How to think on this?

- Answer. Heisenberg group:

- $H = \underbrace{\mathbb{Z}_N \times \mathbb{Z}_N}_v \times \underbrace{\mathbb{Z}_N}_z;$

- $(v, z) \cdot (v', z') = (v + v', z + z' + \frac{1}{2} \begin{vmatrix} v & \\ & v' \end{vmatrix});$

- $(0, 0) \cdot (v, z) = (v, z) \cdot (0, 0) = (v, z);$

- $(v, z) \cdot (-v, -z) = (-v, -z) \cdot (v, z) = (0, 0).$

Heisenberg Rep'n, cont.

- Summary: Heisenberg Rep'n

$$\left\{ \begin{array}{l} \pi : H \rightarrow GL(\mathcal{H}); \\ \pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{\frac{1}{2}\omega\tau + z\}} \cdot M_\omega \circ L_\tau; \\ \pi(h \cdot h') = \pi(h) \circ \pi(h') \text{ — homomorphism.} \end{array} \right.$$

Heisenberg Rep'n, cont.

- Summary: Heisenberg Rep'n

$$\left\{ \begin{array}{l} \pi : H \rightarrow GL(\mathcal{H}); \\ \pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau; \\ \pi(h \cdot h') = \pi(h) \circ \pi(h') \text{ — homomorphism.} \end{array} \right.$$

Definition

A representation of a group H on a complex vector space \mathcal{H} is a homomorphism

$$\begin{aligned} \pi & : H \rightarrow GL(\mathcal{H}), \\ \pi(h \cdot h') & = \pi(h) \circ \pi(h'). \end{aligned}$$

Heisenberg Rep'n, cont.

- Summary: Heisenberg Rep'n

$$\left\{ \begin{array}{l} \pi : H \rightarrow GL(\mathcal{H}); \\ \pi(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau; \\ \pi(h \cdot h') = \pi(h) \circ \pi(h') \text{ — homomorphism.} \end{array} \right.$$

Definition

A representation of a group H on a complex vector space \mathcal{H} is a homomorphism

$$\begin{aligned} \pi & : H \rightarrow GL(\mathcal{H}), \\ \pi(h \cdot h') & = \pi(h) \circ \pi(h'). \end{aligned}$$

- Question. DFT ?

(II) Representation Theory

Definitions

We say that $(\pi_1, H, \mathcal{H}_1)$, $(\pi_2, H, \mathcal{H}_2)$ are equivalent, $\pi_1 \simeq \pi_2$, if

$$\exists \alpha : \mathcal{H}_1 \xrightarrow{\sim} \mathcal{H}_2 \text{ — Intertwiner,}$$

such that for every $h \in H$

$$\begin{array}{ccc} \mathcal{H}_1 & \xrightarrow{\pi_1(h)} & \mathcal{H}_1 \\ \downarrow \alpha & & \downarrow \alpha \\ \mathcal{H}_2 & \xrightarrow{\pi_2(h)} & \mathcal{H}_2 \end{array}$$

i.e., $\alpha \circ \pi_1(h) = \pi_2(h) \circ \alpha$.

(II) Representation Theory

Definitions

We say that $(\pi_1, H, \mathcal{H}_1)$, $(\pi_2, H, \mathcal{H}_2)$ are equivalent, $\pi_1 \simeq \pi_2$, if

$$\exists \alpha : \mathcal{H}_1 \xrightarrow{\sim} \mathcal{H}_2 \text{ — Intertwiner,}$$

such that for every $h \in H$

$$\begin{array}{ccc} \mathcal{H}_1 & \xrightarrow{\pi_1(h)} & \mathcal{H}_1 \\ \downarrow \alpha & & \downarrow \alpha \\ \mathcal{H}_2 & \xrightarrow{\pi_2(h)} & \mathcal{H}_2 \end{array}$$

i.e., $\alpha \circ \pi_1(h) = \pi_2(h) \circ \alpha$.

- Example: *DFT* is an intertwiner!

Rep'n Theory, cont.

- $H = V \times Z = \underbrace{\mathbb{Z}_N}_T \times \underbrace{\mathbb{Z}_N}_W \times \underbrace{\mathbb{Z}_N}_Z$

Rep'n Theory, cont.

- $H = V \times Z = \underbrace{\mathbb{Z}_N}_T \times \underbrace{\mathbb{Z}_N}_W \times \underbrace{\mathbb{Z}_N}_Z$
- Time representation: $\mathcal{H}_T = \mathbb{C}(T)$

$$\begin{cases} \pi_T : H \rightarrow GL(\mathcal{H}_T); \\ \pi_T(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau. \end{cases}$$

Rep'n Theory, cont.

- $H = V \times Z = \underbrace{\mathbb{Z}_N}_T \times \underbrace{\mathbb{Z}_N}_W \times \underbrace{\mathbb{Z}_N}_Z$

- Time representation: $\mathcal{H}_T = \mathbb{C}(T)$

$$\begin{cases} \pi_T : H \rightarrow GL(\mathcal{H}_T); \\ \pi_T(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{\frac{1}{2}\omega\tau + z\}} \cdot M_\omega \circ L_\tau. \end{cases}$$

- Frequency representation: $\mathcal{H}_W = \mathbb{C}(W)$

$$\begin{cases} \pi_W : H \rightarrow GL(\mathcal{H}_W); \\ \pi_W(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{-\frac{1}{2}\omega\tau + z\}} \cdot M_\tau \circ L_{-\omega}. \end{cases}$$

Rep'n Theory, cont.

- $H = V \times Z = \underbrace{\mathbb{Z}_N}_T \times \underbrace{\mathbb{Z}_N}_W \times \underbrace{\mathbb{Z}_N}_Z$

- Time representation: $\mathcal{H}_T = \mathbb{C}(T)$

$$\begin{cases} \pi_T : H \rightarrow GL(\mathcal{H}_T); \\ \pi_T(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ \frac{1}{2} \omega \tau + z \}} \cdot M_\omega \circ L_\tau. \end{cases}$$

- Frequency representation: $\mathcal{H}_W = \mathbb{C}(W)$

$$\begin{cases} \pi_W : H \rightarrow GL(\mathcal{H}_W); \\ \pi_W(\tau, \omega, z) = e^{\frac{2\pi i}{N} \{ -\frac{1}{2} \omega \tau + z \}} \cdot M_\tau \circ L_{-\omega}. \end{cases}$$

- Fact:

$$DFT \circ \pi_T(h) = \pi_W(h) \circ DFT,$$

where

$$DFT[f](w) = \frac{1}{\sqrt{N}} \sum_{t \in \mathbb{Z}_N} f(t) e^{\frac{2\pi i}{N} wt}.$$

(III) FFT Algorithm

- Idea:

$$\begin{array}{ccc} \mathcal{H}_T & \xrightarrow{\text{fast}} & \mathcal{H}^\Lambda \\ \text{slow} \downarrow \text{DFT} & & \parallel \\ \mathcal{H}_W & \xleftarrow{\text{fast}} & \mathcal{H}^\Lambda \end{array}$$

(III) FFT Algorithm

- Idea:

$$\begin{array}{ccc} \mathcal{H}_T & \xrightarrow{\text{fast}} & \mathcal{H}^\Lambda \\ \text{slow} \downarrow \text{DFT} & & \parallel \\ \mathcal{H}_W & \xleftarrow{\text{fast}} & \mathcal{H}^\Lambda \end{array}$$

- More representation theory:

(III) FFT Algorithm

- Idea:

$$\begin{array}{ccc} \mathcal{H}_T & \xrightarrow{\text{fast}} & \mathcal{H}^\Lambda \\ \text{slow} \downarrow \text{DFT} & & \parallel \\ \mathcal{H}_W & \xleftarrow{\text{fast}} & \mathcal{H}^\Lambda \end{array}$$

- More representation theory:

Definition

A rep'n (π, H, \mathcal{H}) is irreducible if

$$\nexists 0 \neq \mathcal{H}' \subsetneq \mathcal{H}$$

such that

$$\pi(h) \cdot \mathcal{H}' \subset \mathcal{H}', \quad \forall h \in H.$$

Theorem (Stone–von Neumann)

If $(\pi_j, H, \mathcal{H}_j)$, $j = 1, 2$, are irreducible representations of the Heisenberg group $H = V \times Z$, with

$$\pi_j(z) = e^{\frac{2\pi i}{N}z} \cdot \text{Id}_{\mathcal{H}_j}, \quad \forall z \in Z,$$

then $\pi_1 \simeq \pi_2$ are equivalent, i.e., $\exists \alpha : \mathcal{H}_1 \xrightarrow{\sim} \mathcal{H}_2$ such that

$$\alpha \circ \pi_1(h) = \pi_2(h) \circ \alpha, \quad \forall h \in H. \quad (1)$$

Theorem (Stone–von Neumann)

If $(\pi_j, H, \mathcal{H}_j)$, $j = 1, 2$, are irreducible representations of the Heisenberg group $H = V \times Z$, with

$$\pi_j(z) = e^{\frac{2\pi i}{N}z} \cdot \text{Id}_{\mathcal{H}_j}, \quad \forall z \in Z,$$

then $\pi_1 \simeq \pi_2$ are equivalent, i.e., $\exists \alpha : \mathcal{H}_1 \xrightarrow{\sim} \mathcal{H}_2$ such that

$$\alpha \circ \pi_1(h) = \pi_2(h) \circ \alpha, \quad \forall h \in H. \quad (1)$$

Theorem (Schur's lemma)

If $\pi_1 \simeq \pi_2$ equivalent irreducible representations, and if α, α' satisfy equation (1), then

$$\alpha = c \cdot \alpha', \quad \text{for some scalar } c.$$

Examples

(1) Time model: $V \supset T = \{(t, 0); t \in \mathbb{Z}_N\}$

$$\pi_T : H \rightarrow GL(\mathcal{H}_T).$$

(2) Frequency model: $V \supset W = \{(0, w); w \in \mathbb{Z}_N\}$

$$\pi_W : H \rightarrow GL(\mathcal{H}_W).$$

Corollary

DFT : $\mathcal{H}_T \rightarrow \mathcal{H}_W$ — unique (up to a scalar) intertwiner between π_T and π_W .

Examples

(3) W -invariant model:

Space:

$$\mathcal{H}^W = \left\{ \begin{array}{l} f : H \rightarrow \mathbb{C}, \\ f(w \cdot h) = f(h), \quad \forall w \in W, h \in H, \\ f(z \cdot h) = e^{\frac{2\pi i}{N}z} \cdot f(h), \quad \forall z \in Z, h \in H. \end{array} \right.$$

Action:

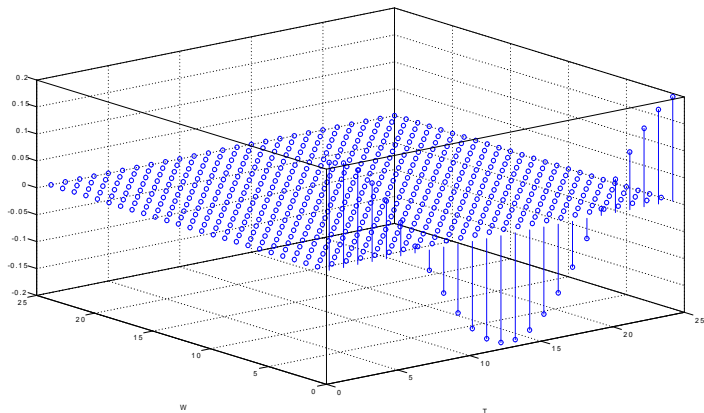
$$\left\{ \begin{array}{l} \pi^W : H \rightarrow GL(\mathcal{H}^W), \\ [\pi^W(h')f](h) = f(h \cdot h'). \end{array} \right.$$

- Remark: We have a natural identification $\mathcal{H}_T = \mathcal{H}^W$, given by

$$f \mapsto \tilde{f}(wtz) = e^{\frac{2\pi i}{N}z} \cdot f(t).$$

Think on W -invariant functions as functions on T

A function $f(t)$ on $T = \{(t, 0, 0); t \in \mathbb{Z}/5^2\}$



Examples

(4) T -invariant model:

Space:

$$\mathcal{H}^T = \begin{cases} g : H \rightarrow \mathbb{C}, \\ g(t \cdot h) = g(h), \quad \forall t \in T, h \in H, \\ g(z \cdot h) = e^{\frac{2\pi i}{N}z} \cdot g(h), \quad \forall z \in Z, h \in H. \end{cases}$$

Action:

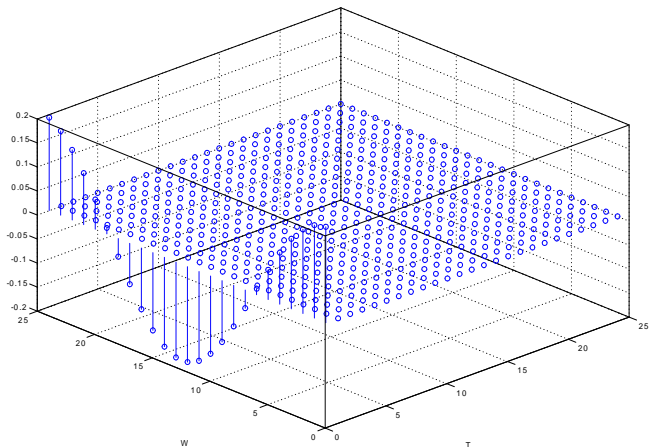
$$\begin{cases} \pi^T : H \rightarrow GL(\mathcal{H}^T), \\ [\pi^T(h')g](h) = g(h \cdot h'). \end{cases}$$

- Remark: We have a natural identification $\mathcal{H}_W = \mathcal{H}^T$, given by

$$g \mapsto \tilde{g}(twz) = e^{\frac{2\pi i}{N}z} \cdot g(w).$$

Think on T-invariant functions as functions on W

A function $g(w)$ on $W = \{(0, w, 0); w \in \mathbb{Z}/5^2\}$



FFT Algorithm: The Arithmetic Model

Examples

(5) Arithmetic model: $N = p^2$

Lagrangian: $V = \mathbb{Z}/p^2 \times \mathbb{Z}/p^2 \supset \Lambda = \{(p \cdot a, p \cdot b)\}$

Space:

$$\mathcal{H}^\Lambda = \begin{cases} F : H \rightarrow \mathbb{C}, \\ F(\lambda \cdot h) = F(h), \quad \forall \lambda \in \Lambda, h \in H, \\ F(z \cdot h) = e^{\frac{2\pi i}{N} z} \cdot F(h), \quad \forall z \in Z, h \in H. \end{cases}$$

Action:

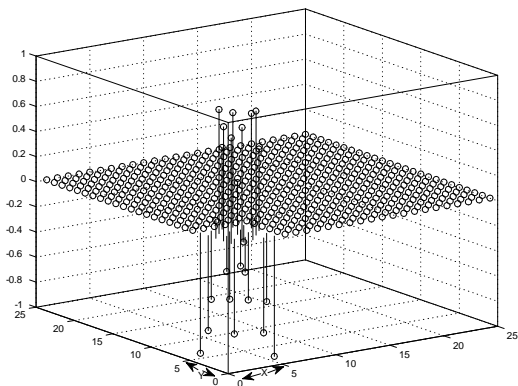
$$\pi^\Lambda : H \rightarrow GL(\mathcal{H}^\Lambda), \quad [\pi^\Lambda(h')F](h) = F(h \cdot h').$$

- Remark: We have a natural identification of \mathcal{H}^Λ with functions $F(x, y)$ on $\{0, \dots, p-1\} \times \{0, \dots, p-1\}$

$$F \mapsto \tilde{F}(\lambda \cdot (x, y, 0) \cdot z) = e^{\frac{2\pi i}{N} z} \cdot F(x, y), \quad \lambda \in \Lambda.$$

Think on Lambda-invariant functions as functions on $\{0, \dots, p-1\} \times \{0, \dots, p-1\}$

A function $F(x, y)$ on $(x, y) \in \{0, \dots, 4\} \times \{0, \dots, 4\}$, $p = 5$,



FFT: The Algorithm

- Algorithm: Case $N = p^2$

$$\begin{array}{ccc} \mathcal{H}_T & \xrightarrow[p^4]{DFT} & \mathcal{H}_W \\ \text{saw} \downarrow & & \uparrow \text{saw} \\ \mathcal{H}^W & & \mathcal{H}^T \\ \Sigma_{\lambda \in \Lambda} f(\lambda h) \downarrow ? & & ? \uparrow \Sigma_{t \in T} F(tw) \\ \mathcal{H}^\Lambda & \underline{\underline{=}} & \mathcal{H}^\Lambda \end{array}$$

FFT: The Algorithm

- Algorithm: Case $N = p^2$

$$\begin{array}{ccc} \mathcal{H}_T & \xrightarrow[p^4]{DFT} & \mathcal{H}_W \\ \text{saw} \downarrow & & \uparrow \text{saw} \\ \mathcal{H}^W & & \mathcal{H}^T \\ \Sigma_{\lambda \in \Lambda} f(\lambda h) \downarrow ? & & ? \uparrow \Sigma_{t \in T} F(th) \\ \mathcal{H}^\Lambda & \text{=====} & \mathcal{H}^\Lambda \end{array}$$

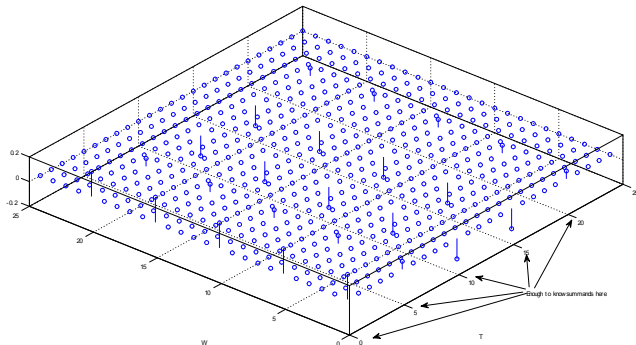
- Observation (**the saving!**): For $f \in \mathcal{H}^W$, if $\lambda, \lambda' \in \Lambda$ differ by an element of W , i.e., $\lambda = w\lambda'$, then

$$f(\lambda h) = f(w\lambda' h) = f(\lambda' h), \quad \forall h \in H.$$

FFT Algorithms: The Summands

So for a fixed h it is enough to know the summands $f(\lambda h)$, in the transform, only for $\lambda \in \Lambda/[\Lambda \cap W] = \{(0, 0), (p, 0), \dots, ((p-1)p, 0)\}$.

- Example for $p = 5$



FFT Algorithm: The Saving

- So we have

$$\sum_{\lambda \in \Lambda} f(\lambda h) = \sum_{\lambda \in \Lambda / [\Lambda \cap W]} f(\lambda h) \cdot \#(\Lambda \cap W).$$

Only $p = p^2 / p = \#(\Lambda / [\Lambda \cap W])$ summands !!!.

FFT Algorithm: The Saving

- So we have

$$\sum_{\lambda \in \Lambda} f(\lambda h) = \sum_{\lambda \in \Lambda / [\Lambda \cap W]} f(\lambda h) \cdot \#(\Lambda \cap W).$$

Only $p = p^2 / p = \#(\Lambda / [\Lambda \cap W])$ summands !!!.

- Complexity of the algorithm

$$\underbrace{p^2}_{\text{values of } h} \cdot \underbrace{p}_{\text{summands}} + \underbrace{p^2 \cdot p}_{\text{second operator}} = p \cdot p^2 \cdot (1 + 1) = \underbrace{p}_{\text{constant}} \cdot N \cdot \log(N).$$

FFT Algorithm: Schur's lemma

Denote the intertwiners by

$$FFT^{\Lambda, W} : \mathcal{H}^W \rightarrow \mathcal{H}^\Lambda, \text{ and } FFT^{T, \Lambda} : \mathcal{H}^\Lambda \rightarrow \mathcal{H}^T.$$

- Why

$$DFT = FFT^{T, \Lambda} \circ FFT^{\Lambda, W} ? \quad (2)$$

FFT Algorithm: Schur's lemma

Denote the intertwiners by

$$FFT^{\Lambda, W} : \mathcal{H}^W \rightarrow \mathcal{H}^\Lambda, \text{ and } FFT^{T, \Lambda} : \mathcal{H}^\Lambda \rightarrow \mathcal{H}^T.$$

- Why

$$DFT = FFT^{T, \Lambda} \circ FFT^{\Lambda, W} ? \quad (2)$$

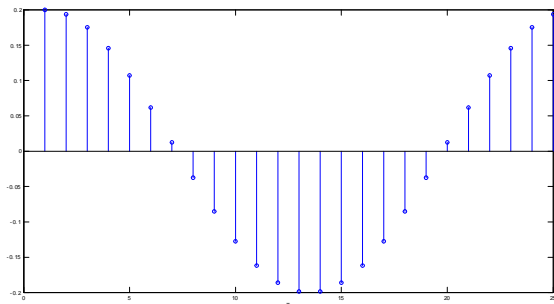
- In fact by Schur's lemma: Both sides of (2) intertwine π^W and π^T , so

$$DFT = c \cdot FFT^{T, \Lambda} \circ FFT^{\Lambda, W}$$

for some scalar c . Easy to compute $c = \frac{1}{p}$.

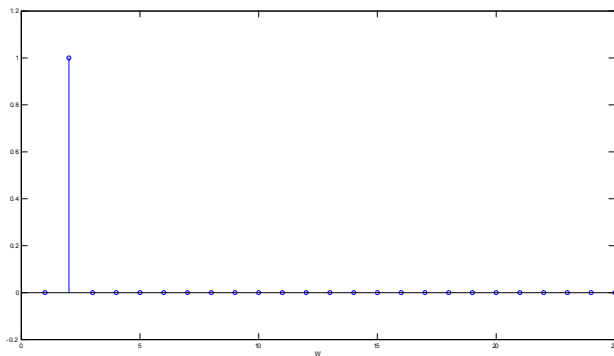
Numerics for $n=25$

Start with the exponent function $f(t) = e^{\frac{2\pi i}{25}t}/5$ on $T = \mathbb{Z}/25$



Numerics for $n=25$: $FFT(f)$

Apply FFT and get $g(w) = FFT[f](w)$ on $W = \mathbb{Z}/25$



Indeed this is $g(w) = \delta_1(w)$.

FFT Algorithm: General

- $N = p^k, k \geq 1$

FFT Algorithm: General

- $N = p^k, k \geq 1$
- $V = \mathbb{Z}/p^k \times \mathbb{Z}/p^k$

FFT Algorithm: General

- $N = p^k$, $k \geq 1$
- $V = \mathbb{Z}/p^k \times \mathbb{Z}/p^k$
- Good Lagrangian sequence ("interpolating" W and T)

$$\Lambda_{\bullet} : \quad \Lambda_j = \{(p^{k-j} \cdot a, p^j \cdot b)\}, \quad j = 0, \dots, k,$$

with LARGE intersections

$$\# \Lambda_j \cap \Lambda_{j+1} = p^{k-1}.$$

FFT Algorithm: General

- $N = p^k, k \geq 1$
- $V = \mathbb{Z}/p^k \times \mathbb{Z}/p^k$
- Good Lagrangian sequence ("interpolating" W and T)

$$\Lambda_\bullet: \quad \Lambda_j = \{(p^{k-j} \cdot a, p^j \cdot b)\}, \quad j = 0, \dots, k,$$

with LARGE intersections

$$\# \Lambda_j \cap \Lambda_{j+1} = p^{k-1}.$$

- FFT Algorithm

$$\begin{aligned} \mathcal{H}^W &\xrightarrow{FFT^{\Lambda_1, W}} \mathcal{H}^{\Lambda_1} \rightarrow \dots \rightarrow \mathcal{H}^{\Lambda_{k-1}} \xrightarrow{FFT^{T, \Lambda_k}} \mathcal{H}^T \\ DFT &= \frac{1}{\sqrt{p^k}} \cdot FFT^{T, \Lambda_{k-1}} \circ \dots \circ FFT^{\Lambda_1, W}, \end{aligned}$$

with complexity

$$\underbrace{p^k}_{\text{values of } h} \cdot \underbrace{\{k\}}_{k \text{ operators}} \cdot \underbrace{\# \Lambda_{j+1} / (\Lambda_j \cap \Lambda_{j+1})}_{p \text{ summands}} = \underbrace{p}_{\text{constant}} \cdot \underbrace{p^k}_N \cdot \underbrace{\log(p^k)}_{\log(N)}.$$

- Conclusion

$\Lambda_{\bullet} \implies$ Cooley–Tukey *FFT*

- Conclusion

$\Lambda_{\bullet} \implies$ Cooley–Tukey *FFT*

- Open Problem

$\Lambda_{\text{new}} \implies$ new *FFT!!!*

THANK YOU



Ronny